

Counting Homomorphisms from Hypergraphs of Bounded Generalised Hypertree Width

A Logical Characterisation

Benjamin Scheidt Nicole Schweikardt
Humboldt-Universität zu Berlin



Abstract

- Investigate the distinguishing power of homomorphisms from hypergraphs of bounded generalised hypertree width.
- We introduce a simple 2-sorted counting logic GC^k that expresses properties over hypergraphs.
- We show: Two hypergraphs are distinguishable by the number of homomorphisms from a hypergraph of generalised hypertree width at most k if and only if they can be distinguished by some sentence of the logic GC^k .

Side Result

For hypergraphs G and H it holds that

$$\begin{aligned} \text{Hom}(\text{GHW}_k, G) &= \text{Hom}(\text{GHW}_k, H) \\ &\iff \\ \text{Hom}(\text{EHW}_k, G) &= \text{Hom}(\text{EHW}_k, H) \end{aligned}$$

“For homomorphism indistinguishability it is enough to consider the more restrictive entangled hypertree width.”

The Logic GC^k

Variables

- Countably many **red** variables v_1, v_2, v_3, \dots representing vertices.
- k **blue** variables e_1, e_2, \dots, e_k representing hyperedges.

Atomic formulas

- $E(e_i, v_j), v_i = v_j, e_i = e_j$

Negation, Conjunction and Disjunction

- $\neg\varphi, (\varphi \wedge \psi), (\varphi \vee \psi)$

Guarded Quantification

- $\exists^{\geq n}(v_{i_1}, \dots, v_{i_\ell}).(\Delta_g \wedge \varphi)$
- $\exists^{\geq n}(e_{i_1}, \dots, e_{i_\ell}).(\Delta_g \wedge \varphi)$

where Δ_g has to guard all free red variables in φ .

Main Result

For two hypergraphs G and H and any natural number k it holds that

$$G \equiv_{GC^k} H \iff \text{Hom}(\text{GHW}_k, G) = \text{Hom}(\text{GHW}_k, H).$$

“ G and H are distinguishable by the number of homomorphisms from a hypergraph F of $\text{ghw} \leq k$ iff they are distinguishable by a sentence of the logic GC^k .”

“All hyperedges with ≥ 3 vertices are disjoint.”

$$\begin{aligned} \psi_1 &:= \exists^{\geq 1}(v_1).(E(e_1, v_1) \wedge E(e_2, v_1)) & \psi_2 &:= \bigwedge_{i \in \{1,2\}} \exists^{\geq 3}(v_1).(E(e_i, v_1)) \\ \varphi &:= \neg \exists^{\geq 1}(e_1, e_2).(\neg e_1 = e_2 \wedge \psi_1 \wedge \psi_2) \end{aligned}$$

Proof Overview

$$\begin{aligned} \text{Hom}(\text{GHW}_k, G) &= \text{Hom}(\text{GHW}_k, H) && [\text{Böker 2019}\ddagger] \\ &\iff && \\ \text{Hom}(\text{IGHW}_k, I_G) &= \text{Hom}(\text{IGHW}_k, I_H) && [\text{Thm. 4.1}\ddagger] \\ &\iff && \\ \text{Hom}(\text{IEHW}_k, I_G) &= \text{Hom}(\text{IEHW}_k, I_H) && [\text{Thm. 7.12}\ddagger] \\ &\iff && \\ \text{Hom}(\text{GLI}_k, I_G) &= \text{Hom}(\text{GLI}_k, I_H) && [\text{Thm. 6.1}\ddagger] \\ &\iff && \\ G &\equiv_{\text{RGC}^k} H && \\ &\iff && [\text{Thm. TBA}] \\ G &\equiv_{GC^k} H \end{aligned}$$

†: arXiv:2303.10980v2 [cs.LO] – use the QR code above.

‡: Böker, J. *Color Refinement, Homomorphisms, and Hypergraphs*. WG 2019.

Hypertree Decompositions

Generalised Hypertree Decomp's

- Bags (set of vertices) and covers (set of hyperedges) associated with each node of the tree.
- F.a. vertices v : the bags containing v must form a connected subtree
- Every bag has to be covered by the set of hyperedges.

Entangled Hypertree Decomp's

- Every EHD is a GHD.
- F.a. hyperedges e : the bags covered by e must form a connected subtree.
- Every bag has to be *precisely* covered by the set of hyperedges.

Example of a GHD that is not an EHD

